

La. 2.2.5. If  $X$  and  $Y$  are discrete i.i.d. r.v. with entropy  $H(X)$ . Then

$$P(X=Y) \geq 2^{-H(X)} \quad \square$$

Proof. Let  $p(x)$  denote the pmf of  $X$  and  $Y$ .

w.l.o.g. choose log-base = 2.

$f(t) = 2^t$  is a convex fct. Hence,

$$\begin{aligned} 2^{-H(X)} &= 2^{E(\log p(X))} \\ &\leq E(2^{\log p(X)}) \quad (\text{by Jensen inequ.}) \\ &= \sum_x p(x) 2^{\log p(x)} \\ &= \sum_x p^2(x) = P(X=Y). \quad \square \end{aligned}$$

Is this inequ. sharp?

a) Let  $X, Y$  i.i.d.  $\sim U(1, \dots, m)$  (uniformly distributed)

$$\left. \begin{aligned} H(X) &= \log m, \quad 2^{-H(X)} = \frac{1}{m} \\ P(X=Y) &= \sum_{i=1}^m \frac{1}{m^2} = \frac{1}{m} \end{aligned} \right\} =$$

b)  $X, Y$  i.i.d. with only one mass point.

$$\left. \begin{aligned} H(X) &= 0, \quad 2^{-H(X)} = 1 \\ P(X=Y) &= 1 \end{aligned} \right\} =$$

Inequality is sharp at least for two cases.

## 2.3. Information Measures for Random Sequences

Consider sequences of r.v.  $X_1, X_2, X_3, \dots$   
denoted as  $X = \{X_n\}_{n \in \mathbb{N}}$ .

Naive approach: define the entropy of  $X$

$$H(X) = \lim_{n \rightarrow \infty} H(X_1, X_2, \dots, X_n).$$

In most cases this limit will be infinite.  
Instead consider the entropy rate.

Def. 2.3.1. Let  $X = \{X_n\}_{n \in \mathbb{N}}$  be a sequence of discrete r.v.

$$H_\infty(X) = \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, \dots, X_n)$$

is called the entropy rate of  $X$ , provided the limit exists.

$H_\infty(X)$  is the average uncertainty per symbol.

Def. 2.3.2. Let  $X = \{X_n\}_{n \in \mathbb{N}}$ ,  $Y = \{Y_n\}_{n \in \mathbb{N}}$  sequ. of discr. r.v.

$$I_\infty(X, Y) = \lim_{n \rightarrow \infty} \frac{1}{n} I(X_1, \dots, X_n; Y_1, \dots, Y_n)$$

is called mutual information rate of  $X$  and  $Y$ .

Example 2.3.2

a) Let  $X = \{X_n\}_{n \in \mathbb{N}}$  be i.i.d. r.v. with  $H(X_i) < \infty$ .

Then

$$\begin{aligned} H_\infty(X) &= \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, \dots, X_n) \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n H(X_i) = H(X_1) \end{aligned}$$

b) Let  $\{Z_n\}_{n \in \mathbb{N}} = \{(X_n, Y_n)\}_{n \in \mathbb{N}}$  be iid sequence with  $I(X_k; Y_k) < \infty$ . Then

$$I_\infty(X; Y) = \lim_{n \rightarrow \infty} \frac{1}{n} I(X_1, \dots, X_n; Y_1, \dots, Y_n)$$

$$\stackrel{(\text{Ex.})}{=} \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{k=1}^n I(X_k; Y_k)$$

$$= I(X_1; Y_1) \quad \square$$

Going further than iid. sequences.

Def. 2.3.3. A sequence of r.v.s  $X = \{X_n\}_{n \in \mathbb{N}}$  is called (strongly) stationary if

$$P(X_{i_1}, \dots, X_{i_k}) = P(X_{i_1+t}, \dots, X_{i_k+t})$$

for  $\forall 1 \leq i_1 < \dots < i_k, t \in \mathbb{N}$ .

- The joint distribution of any finite selection of r.v.s from  $\{X_n\}_{n \in \mathbb{N}}$  is invariant w.r.t. time shift.
- An equivalent condition for discrete r.v. with support  $\mathcal{X}$  is as follows

$$P(X_1 = s_1, \dots, X_n = s_n) = P(X_{1+t} = s_1, \dots, X_{n+t} = s_n)$$

for all  $s_1, \dots, s_n \in \mathcal{X}, n \in \mathbb{N}, t \in \mathbb{N}$ .

- For stationary sequences all marginal distributions  $P_{X_t}$  are the same

Theorem 2.3.4. Let  $X = \{X_n\}_{n \in \mathbb{N}}$  be a stationary seq. Then

- a)  $H(X_n | X_{1, \dots, n-1})$  is monotonically decreasing.
- b)  $H(X_n | X_{1, \dots, n-1}) \leq \frac{1}{n} H(X_{1, \dots, n})$
- c)  $\frac{1}{n} H(X_{1, \dots, n})$  is monotonically decreasing.
- d)  $\lim_{n \rightarrow \infty} H(X_n | X_{1, \dots, n-1}) = \lim_{n \rightarrow \infty} \frac{1}{n} H(X_{1, \dots, n}) = H_\infty(X)$

Because of a) and c) both limits exist.  $\square$

Proof.

$$a) \quad H(X_n | X_{1, \dots, n-1}) \stackrel{\substack{\uparrow \\ \text{Th. 2.1.8 d)}}}{\leq} H(X_n | X_{2, \dots, n-1}) \stackrel{\substack{\uparrow \\ \text{stationary}}}{\leq} H(X_{n-1} | X_{1, \dots, n-2})$$

Hence  $H(X_n | X_{1, \dots, n-1})$  is mon. decreasing.

The first limit in d) exists since 0 is a lower bound.

b) By Th. 2.1.5 (chain rule)

$$\frac{1}{n} H(X_{1, \dots, n}) = \frac{1}{n} (H(X_1) + H(X_2 | X_1) + \dots + H(X_n | X_{1, \dots, n-1}))$$

$$\stackrel{a)}{\geq} H(X_n | X_{1, \dots, n-1})$$

$$c) \quad H(X_{1, \dots, n}) = H(X_{1, \dots, n-1}) + H(X_n | X_{1, \dots, n-1})$$

$$\stackrel{b)}{\leq} H(X_{1, \dots, n-1}) + \frac{1}{n} H(X_{1, \dots, n})$$

Hence,

$$\frac{n-1}{n} H(X_{1, \dots, n}) \leq H(X_{1, \dots, n-1})$$

$$\Leftrightarrow \frac{1}{n} H(X_{1, \dots, n}) \leq \frac{1}{n-1} H(X_{1, \dots, n-1})$$

which proves monotonicity.

$$\begin{aligned}
 d) \quad & \frac{1}{n+k} H(X_1, \dots, X_{n+k}) \\
 &= \frac{1}{n+k} \left[ H(X_{n+k} | X_1, \dots, X_{n+k-1}) + \dots + H(X_{n+1} | X_1, \dots, X_n) \right. \\
 &\quad \left. + H(X_n | X_1, \dots, X_{n-1}) + H(X_1, \dots, X_{n-1}) \right] \\
 &\leq \underbrace{\frac{1}{n+k} H(X_1, \dots, X_{n-1})}_{\rightarrow 0 \text{ (} k \rightarrow \infty)} + \underbrace{\frac{k+1}{n+k} H(X_n | X_1, \dots, X_{n-1})}_{\rightarrow H(X_n | X_1, \dots, X_{n-1}) \text{ (} k \rightarrow \infty)}
 \end{aligned}$$

Now fix  $n$ , set  $l = n+k$

Using b)

$$\begin{aligned}
 \lim_{n \rightarrow \infty} H(X_n | X_1, \dots, X_{n-1}) &\leq \lim_{n \rightarrow \infty} \frac{1}{n} H(X_1, \dots, X_n) \\
 &\leq \lim_{n \rightarrow \infty} H(X_n | X_1, \dots, X_{n-1})
 \end{aligned}$$

so that the limits ~~are~~ exist and are equal.  $\square$

Example 2.3.5. (English text)

Frequencies of single characters

A	B	...	E	...	Z
0.082	0.015		0.127		0.001

Frequencies of digrams (in %)

AN	TH	...	TI
1.81	3.21		1.28

Estimated entropy rates (log-base = 2)

$n$	1	2	3	...
$\frac{1}{n} H(X_{1..n})$	4.14	3.56	3.3	

From experiments it is estimated (log-base = 2)

$$1 < H_{\infty}(X) \leq 1.5 !$$

# English Letter Frequencies

The frequencies from this page are generated from around 4.5 billion characters of English text, sourced from [Wortschatz](#). The text files containing the counts can be used with `ngram_score.py` and used for breaking ciphers, see [this page](#) for details. If you want to compute the letter frequencies of your own piece of text you can use [this page](#).

English single letter frequencies are as follows (in percent %):

A : 8.55      K : 0.81      U : 2.68  
 B : 1.60      L : 4.21      V : 1.06  
 C : 3.16      M : 2.53      W : 1.83  
 D : 3.87      N : 7.17      X : 0.19  
 E : 12.10    O : 7.47      Y : 1.72  
 F : 2.18      P : 2.07      Z : 0.11  
 G : 2.09      Q : 0.10  
 H : 4.96      R : 6.33  
 I : 7.33      S : 6.73  
 J : 0.22      T : 8.94

Bigrams. The top 30 are the following (in percent %):

TH : 2.71      EN : 1.13      NG : 0.89  
 HE : 2.33      AT : 1.12      AL : 0.88  
 IN : 2.03      ED : 1.08      IT : 0.88  
 ER : 1.78      ND : 1.07      AS : 0.87  
 AN : 1.61      TO : 1.07      IS : 0.86  
 RE : 1.41      OR : 1.06      HA : 0.83  
 ES : 1.32      EA : 1.00      ET : 0.76  
 ON : 1.32      TI : 0.99      SE : 0.73  
 ST : 1.25      AR : 0.98      OU : 0.72  
 NT : 1.17      TE : 0.98      OF : 0.71

Trigrams. The top 30 are the following (in percent %):

THE : 1.81      ERE : 0.31      HES : 0.24  
 AND : 0.73      TIO : 0.31      VER : 0.24  
 ING : 0.72      TER : 0.30      HIS : 0.24  
 ENT : 0.42      EST : 0.28      OFT : 0.22  
 ION : 0.42      ERS : 0.28      ITH : 0.21  
 HER : 0.36      ATI : 0.26      FTH : 0.21  
 FOR : 0.34      HAT : 0.26      STH : 0.21  
 THA : 0.33      ATE : 0.25      OTH : 0.21  
 NTH : 0.33      ALL : 0.25      RES : 0.21  
 INT : 0.32      ETH : 0.24      ONT : 0.20

Quadgrams. The top 30 are the following (in percent %):

TION : 0.31      OTHE : 0.16      THEM : 0.12  
 NTHE : 0.27      TTHE : 0.16      RTHE : 0.12  
 THER : 0.24      DTHE : 0.15      TREP : 0.11  
 THAT : 0.21      INGT : 0.15      FROM : 0.10  
 OFTH : 0.19      ETHE : 0.15      THIS : 0.10  
 FTHE : 0.19      SAND : 0.14      TING : 0.10  
 THES : 0.18      STHE : 0.14      THEI : 0.10  
 WITH : 0.18      HERE : 0.13      NGTH : 0.10  
 INTH : 0.17      THEC : 0.13      IONS : 0.10  
 ATIO : 0.17      MENT : 0.12      ANDT : 0.10

Quintigrams. The top 30 are the following (in percent %):

OFTHE : 0.18      ANDTH : 0.07      CTION : 0.05  
 ATION : 0.17      NDTHE : 0.07      WHICH : 0.05  
 INTHE : 0.16      ONTHE : 0.07      THESE : 0.05  
 THERE : 0.09      EDTHE : 0.06      AFTER : 0.05  
 INGTH : 0.09      THEIR : 0.06      EOFTH : 0.05  
 TOTHE : 0.08      TIONA : 0.06      ABOUT : 0.04  
 NGTHE : 0.08      ORTHE : 0.06      ERTHE : 0.04  
 OTHER : 0.07      FORTH : 0.06      IONAL : 0.04  
 ATTHE : 0.07      INGTO : 0.06      FIRST : 0.04  
 TIONS : 0.07      THECO : 0.05      WOULD : 0.04