Lehrstuhl für Theoretische Informationstechnik

Homework 6 in Cryptography I Prof. Dr. Rudolf Mathar, Michael Naehrig 26.11.2007

Exercise 16. Let M be a block of bits of length 64 and K be a block of bits of length 56. Let DES(M, K) denote the encryption of M with key K using the DES cryptosystem. Show that

$$DES(M, K) = \overline{DES(\overline{M}, \overline{K})},$$

where $\bar{\cdot}$ denotes the bitwise complement.

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Exercise 17. Consider the following cryptosystem. Messages are bit sequences of arbitrary length, i.e. character sequences over the alphabet $\mathbb{F}_2 = \{0, 1\}$. Let the message be $m = m_1 m_2 \dots m_l$. Keys are also bit sequences $k = k_1 k_2 \dots k_n$ of fixed length n. Usually we have n < l. Now a key stream $z = z_1 z_2 \dots z_l$ is generated in a recursive manner depending on the key:

$$z_i = k_i, \quad 1 \le i \le n,$$

 $z_i = \sum_{j=1}^n s_j z_{i-j} \pmod{2}, \quad n < i \le l.$

Here s_1, \ldots, s_n are fixed bits which are given in advance. We encrypt $c_i := m_i \oplus z_i$ for $1 \le i \le l$.

- (a) How does decryption work for this cryptosystem? Why should k = 00...0 not be chosen as the key?
- (b) Encrypt the message m = 1011000101001101000 with n = 4, $s_2 = s_3 = 0$, $s_1 = s_4 = 1$ using the key k = 0110. The key stream is periodic. How long is its period?

This method for generating a key stream is called linear feedback shift register (LFSR).

Exercise 18. Consider the finite field \mathbb{F}_4 from Exercise 12. Construct an extension field \mathbb{F}_{16} of \mathbb{F}_4 with 16 elements and describe your approach. Hint: Start with the polynomial ring $\mathbb{F}_4[U]$.